

# iFM & ABZ 2012 Tutorial

## Safety, Dependability and Performance Analysis of Extended AADL Models

### Part 2: System Modeling Using AADL



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European Space Research and Technology Centre



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- 1 Formal Semantics of Nominal Specifications
- 2 Model Extension
- 3 References & Ongoing Activities

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# The Battery Revisited

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;
end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged: activation mode
      while energy'=-2.0 and energy>=20.0;
    depleted: mode
      while energy'=-3.0 and energy>=0.0;
  transitions
    charged -[then voltage:=energy/50.0+4.0]-> charged;
    charged -[empty when energy<=20.0]-> depleted;
    depleted -[then voltage:=energy/50.0+4.0]->
      depleted;
end Battery.Imp;
```

## Definition (Event-data automaton)

An **event-data automaton (EDA)** is a tuple

$$\mathcal{A} = (M, X, V, \iota, E, \rightarrow)$$

with

- $M$  finite set of **modes**
- $X = IX \uplus OX \uplus LX$  finite set of **input/output/local variables**
- $V := \{v \mid v : X \rightarrow \dots\}$  **valuations**
- $\iota : M \rightarrow (V \rightarrow \mathbb{B})$  **mode invariants**
- $E = IE \uplus OE$  finite set of **input/output events**
- $\rightarrow \subseteq M \times \underbrace{(E \cup \{\tau\})}_{\text{trigger}} \times \underbrace{(V \rightarrow \mathbb{B})}_{\text{guard}} \times \underbrace{(V \rightarrow V)}_{\text{effect}} \times M$  **transition relation**

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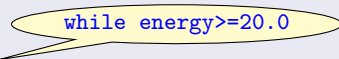
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empty

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when energy ≤ 20.0

empty

then voltage := ...

# Operational Semantics of EDAs

- **States:**  $M \times V$
- **Transitions:** timed or internal or synchronized

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## Example (Battery)

$\langle \text{mode} = \text{charged}, \text{energy} = 100.0, \text{voltage} = 6.0 \rangle$

# Operational Semantics of EDAs

- States:  $M \times V$
- Transitions: **timed** or internal or synchronized

## Example (Battery)

$\langle \text{mode} = \text{charged}, \text{energy} = 100.0, \text{voltage} = 6.0 \rangle$

$\downarrow 30.0$

$\langle \text{mode} = \text{charged}, \text{energy} = 40.0, \text{voltage} = 6.0 \rangle$

# Operational Semantics of EDAs

- States:  $M \times V$
- Transitions: timed or **internal** or synchronized

## Example (Battery)

```
⟨mode = charged, energy = 100.0, voltage = 6.0⟩  
    ↓ 30.0  
⟨mode = charged, energy = 40.0, voltage = 6.0⟩  
    ↓  $\tau$ ⟨voltage := ...⟩  
⟨mode = charged, energy = 40.0, voltage = 4.8⟩
```

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- Transitions: **timed** or internal or synchronized

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⟨mode = charged, energy = 100.0, voltage = 6.0⟩  
    ↓ 30.0  
⟨mode = charged, energy = 40.0, voltage = 6.0⟩  
    ↓  $\tau$ ⟨voltage:=...⟩  
⟨mode = charged, energy = 40.0, voltage = 4.8⟩  
    ↓ 10.0  
⟨mode = charged, energy = 20.0, voltage = 4.8⟩
```

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⟨mode = charged, energy = 40.0, voltage = 4.8⟩  
    ↓ 10.0  
⟨mode = charged, energy = 20.0, voltage = 4.8⟩  
    ↓  $\tau$ ⟨voltage:=...⟩  
⟨mode = charged, energy = 20.0, voltage = 4.4⟩
```



# Operational Semantics of EDAs

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    ↓ τ⟨voltage:=...⟩  
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    ↓ 10.0  
⟨mode = charged, energy = 20.0, voltage = 4.8⟩  
    ↓ τ⟨voltage:=...⟩  
⟨mode = charged, energy = 20.0, voltage = 4.4⟩  
    ↓ empty  
⟨mode = depleted, energy = 20.0, voltage = 4.4⟩
```

# Operational Semantics of EDAs

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## Example (Battery)

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     $\downarrow 30.0$   
 $\langle \text{mode} = \text{charged}, \text{energy} = 40.0, \text{voltage} = 6.0 \rangle$   
     $\downarrow \tau \langle \text{voltage} := \dots \rangle$   
 $\langle \text{mode} = \text{charged}, \text{energy} = 40.0, \text{voltage} = 4.8 \rangle$   
     $\downarrow 10.0$   
 $\langle \text{mode} = \text{charged}, \text{energy} = 20.0, \text{voltage} = 4.8 \rangle$   
     $\downarrow \tau \langle \text{voltage} := \dots \rangle$   
 $\langle \text{mode} = \text{charged}, \text{energy} = 20.0, \text{voltage} = 4.4 \rangle$   
     $\downarrow \text{empty}$   
 $\langle \text{mode} = \text{depleted}, \text{energy} = 20.0, \text{voltage} = 4.4 \rangle$   
     $\downarrow \dots$

# The Power System Revisited

```
system Power
  features
    voltage: out data port real;
end Power;

system implementation Power.Imp
  subcomponents
    batt1: device Battery.Imp in modes (primary);
    batt2: device Battery.Imp in modes (backup);
  connections
    data port batt1.voltage -> voltage
      in modes (primary);
    data port batt2.voltage -> voltage
      in modes (backup);
  modes
    primary: initial mode;
    backup: mode;
  transitions
    primary -[batt1.empty]-> backup;
    backup -[batt2.empty]-> primary;
end Power.Imp;
```

# Networks of Event-Data Automata

## Dynamic reconfiguration

⇒ component activity and port connections **mode dependent**

### Definition (Networks of Event-Data Automata)

A **network of event-data automata (NEDA)** is a tuple

$$\mathfrak{N} = ((\mathfrak{A}_i)_{i \in [n]}, \alpha, EC, DC)$$

with  $n \geq 1$ ,  $[n] := \{1, \dots, n\}$ , and

- each  $\mathfrak{A}_i$  an **EDA**  $\mathfrak{A}_i = (M_i, X_i, V_i, \iota_i, E_i, \rightarrow_i)$
- $M := \prod_{i=1}^n M_i$  set of **global modes**
- $\alpha : M \rightarrow 2^{[n]}$  **activation mapping**
- $EC : M \rightarrow (\{i.e \mid i \in [n], e \in E_i\})^2$  **event connection mapping**
- $DC : M \rightarrow (\{i.x \mid i \in [n], x \in X_i\})^2$  **data connection mapping**

# The Activation Mapping

## Definition (Activation mapping)

$\alpha$  determined by **subcomponent declarations**:

- Root component always **active**
- Component  $c$  **active** and in mode  $m$ ,  
subcomponent  $c'$  of  $c$  activated in  $m$   
 $\Rightarrow c'$  **active**

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 $\Rightarrow c'$  active

## Example (Power system)

For  $\text{Power/Battery1/Battery2}$  ( $m_1, m_2 \in \{\text{charged, depleted}\}$ ):

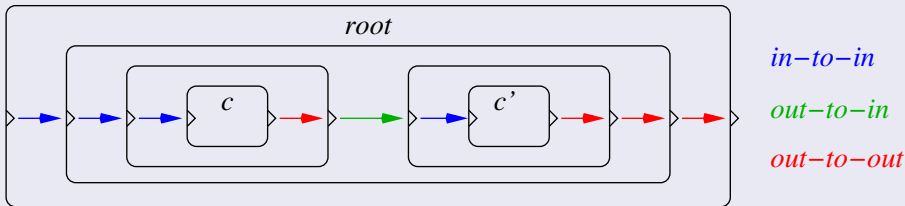
$$\alpha(\text{primary}, m_1, m_2) = \{\text{Power}, \text{Battery1}\}$$

$$\alpha(\text{backup}, m_1, m_2) = \{\text{Power}, \text{Battery2}\}$$

# The Connection Mappings

## Definition (Event/data connection mappings)

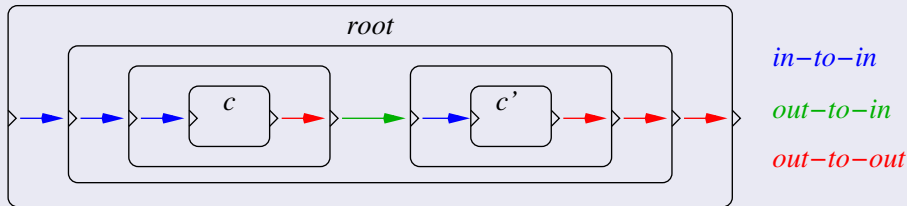
*EC/DC* determined by following all end-to-end **chains of port connections**:



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*EC/DC* determined by following all end-to-end **chains of port connections**:



## Example (Power system)

For *Power/Battery1/Battery2* ( $m_1, m_2 \in \{\text{charged, depleted}\}$ ):

$EC(\text{primary}, m_1, m_2) = \{\text{Battery1.empty} \rightarrow \text{Power.empty}\}$

$EC(\text{backup}, m_1, m_2) = \{\text{Battery2.empty} \rightarrow \text{Power.empty}\}$

$DC(\text{primary}, m_1, m_2) = \{\text{Battery1.voltage} \rightarrow \text{Power.voltage}\}$

$DC(\text{backup}, m_1, m_2) = \{\text{Battery2.voltage} \rightarrow \text{Power.voltage}\}$



# Operational Semantics of Networks of EDAs

- **States**  $:= (M_1 \times V_1) \times \dots \times (M_n \times V_n)$
- **Transitions** determined by active EDAs:
  - 1 Perform local transitions:
    - timed local transition in all EDAs or
    - internal transition in EDA or
    - multi-way event communication from EDA to  $\geq 1$  connected EDAs
  - 2 Initialize (re-)activated subcomponents
  - 3 Establish consistency w.r.t. *DC* (copy source  $\rightarrow$  target data port)

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## Example (Power system)

$\langle m = \underline{\text{primary}}, v = 6.0 \rangle \mid \langle m = \underline{\text{charged}}, e = 100.0, v = 6.0 \rangle \mid \langle m = \text{charged}, e = 100.0, v = 6.0 \rangle$

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 $\quad \downarrow \tau \langle \text{voltage} := \dots \rangle$   
 $\langle m = \underline{\text{primary}}, v = 4.4 \rangle \mid \langle m = \underline{\text{charged}}, e = 20.0, v = 4.4 \rangle \mid \langle m = \text{charged}, e = 100.0, v = 6.0 \rangle$

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 $\langle m = \underline{\text{backup}}, v = 6.0 \rangle \mid \langle m = \text{depleted}, e = 20.0, v = 4.4 \rangle \mid \langle m = \underline{\text{charged}}, e = 100.0, v = 6.0 \rangle$

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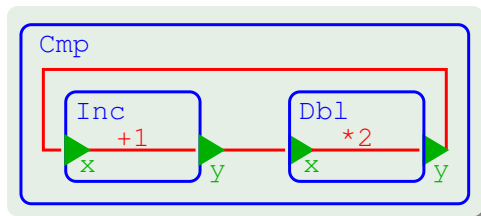
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 $\quad \downarrow \dots$

# AADL Specifications and Attribute Grammars

**Problem:** data port dependencies can be **cyclic**

⇒ data port values possibly **inconsistent**

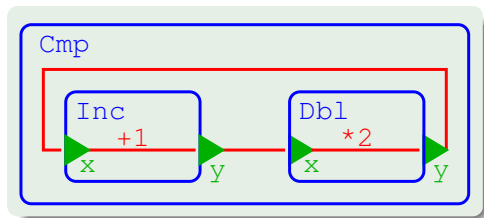




# AADL Specifications and Attribute Grammars

**Problem:** data port dependencies can be **cyclic**

⇒ data port values possibly **inconsistent**

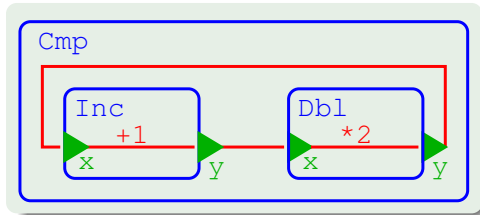


$$\Rightarrow \text{Inc.x} = 2 \cdot (\text{Inc.x} + 1) \quad \text{⚡}$$

# AADL Specifications and Attribute Grammars

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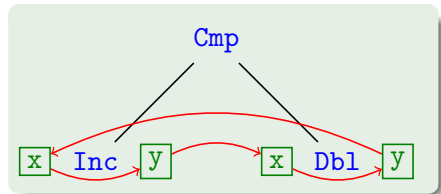
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$$\Rightarrow \text{Inc.x} = 2 \cdot (\text{Inc.x} + 1) \quad \text{⚡}$$

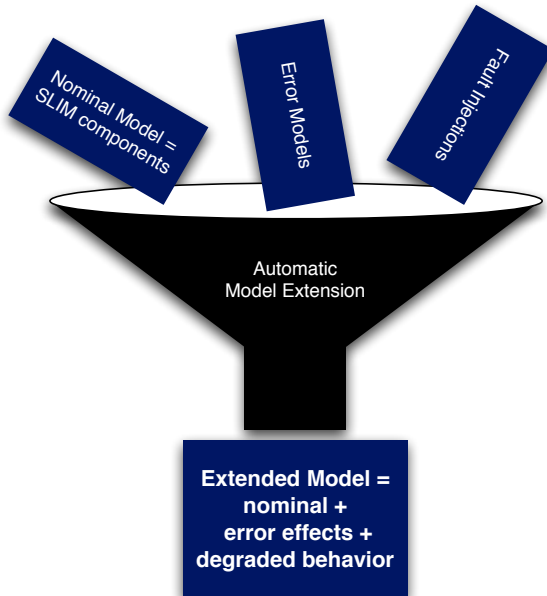
**Equivalent:** **circularity of attribute grammars**

AADL	AGs
Configuration	Derivation tree
Mode	Production
Component	Symbol
Data port	Attribute
Connection	Dependency



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# Integrating Erroneous and Nominal Behavior



# The Battery Error Model Revisited

```
error model BatteryFailure
  features
    ok: initial state;
    dead: error state;
    batteryDied: out error propagation;
end BatteryFailure;

error model implementation BatteryFailure.Imp
  events
    fault: error event occurrence poisson 0.01;
  transitions
    ok -[fault]-> dead;
    dead -[batteryDied]-> dead;
end BatteryFailure.Imp;
```

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  transitions
    ok -[fault]-> dead;
    dead -[batteryDied]-> dead;
end BatteryFailure.Imp;
```

## Fault injection

In error state **dead**, **voltage:=0**

Nominal model + error model + fault injections = **extended model**

- **Modes** are pairs of nominal modes and error model states
  - **Starting mode** = (original starting mode, starting error state)
- **Event ports**  $+=$  error propagations
- **Event port connections**  $+=$  propagation port connections
- **Transition relation**  $:=$  all possible **interleavings** and **interactions** between nominal and error model, taking failure effects into account
- Other elements (e.g., mode invariants) are unaffected

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## Probabilistic error transitions

As an error model has probabilistic transitions, our semantical model has to be equipped with such transitions.

This yields **interactive Markov chains**  $:=$  LTS + Markov chains.



# Battery Component

Nominal specification:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;

end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged: activation mode while ...;
    depleted: mode while ...;
  transitions
    charged -[then voltage:=...]-> charged;
    charged -[empty when energy<=20.0]-> depleted;
    depleted -[then voltage:=...]-> depleted;

end Battery.Imp;
```

# Battery Component **After Model Extension**

**Product construction** for modes:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;

end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged#ok: activation mode while ...;
    depleted#ok, charged#dead, depleted#dead: mode while ...;
  transitions
    charged -[then voltage:=...]-> charged;
    charged -[empty when energy<=20.0]-> depleted;
    depleted -[then voltage:=...]-> depleted;

end Battery.Imp;
```

# Battery Component After Model Extension

Integrate nominal transitions:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;

end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged#ok: activation mode while ...;
    depleted#ok, charged#dead, depleted#dead: mode while ...;
  transitions
    charged#ok -[then voltage:=...]-> charged#ok;
    charged#ok -[empty when energy<=20.0]-> depleted#ok;
    depleted#ok -[then voltage:=...]-> depleted#ok;

end Battery.Imp;
```

# Battery Component After Model Extension

Fault injection:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;

end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged#ok: activation mode while ...;
    depleted#ok, charged#dead, depleted#dead: mode while ...;
  transitions
    charged#ok -[then voltage:=...]-> charged#ok;
    charged#ok -[empty when energy<=20.0]-> depleted#ok;
    depleted#ok -[then voltage:=...]-> depleted#ok;
    charged#ok -[prob 0.01 then voltage:=0]-> charged#dead;
    depleted#ok -[prob 0.01 then voltage:=0]-> depleted#dead;

end Battery.Imp;
```

# Battery Component After Model Extension

Nominal transitions with **fault effects**:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;

end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged#ok: activation mode while ...;
    depleted#ok, charged#dead, depleted#dead: mode while ...;
  transitions
    charged#ok -[then voltage:=...] -> charged#ok;
    charged#ok -[empty when energy<=20.0] -> depleted#ok;
    depleted#ok -[then voltage:=...] -> depleted#ok;
    charged#ok -[prob 0.01 then voltage:=0] -> charged#dead;
    depleted#ok -[prob 0.01 then voltage:=0] -> depleted#dead;
    charged#dead -[then voltage:=0] -> charged#dead;
    charged#dead -[empty when energy<=20.0] -> depleted#dead;
    depleted#dead -[then voltage:=0] -> depleted#dead;

end Battery.Imp;
```

# Battery Component After Model Extension

Add error propagations:

```
device Battery
  features
    empty: out event port;
    voltage: out data port real default 6.0;
    batteryDied: out event port;
end Battery;

device implementation Battery.Imp
  subcomponents
    energy: data continuous default 100.0;
  modes
    charged#ok: activation mode while ...;
    depleted#ok, charged#dead, depleted#dead: mode while ...;
  transitions
    charged#ok -[then voltage:=...]-> charged#ok;
    charged#ok -[empty when energy<=20.0]-> depleted#ok;
    depleted#ok -[then voltage:=...]-> depleted#ok;
    charged#ok -[prob 0.01 then voltage:=0]-> charged#dead;
    depleted#ok -[prob 0.01 then voltage:=0]-> depleted#dead;
    charged#dead -[then voltage:=0]-> charged#dead;
    charged#dead -[empty when energy<=20.0]-> depleted#dead;
    depleted#dead -[then voltage:=0]-> depleted#dead;
    depleted#dead -[batteryDied]-> depleted#dead;
end Battery.Imp;
```

- 1 Formal Semantics of Nominal Specifications
- 2 Model Extension
- 3 References & Ongoing Activities

## References

- Our AADL variant (Bozzano et. al, [MEMOCODE 2009](#))
- AADL formal semantics (Bozzano et. al, [Computer J. 2010](#))
- Relation to attribute grammars (Noll, [FACS 2011](#))



## References

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## Ongoing activities

- Contribution to AADL standardization
- Security aspects in AADL